

GANPAT UNIVERSITY

B. Tech. Semester VI Computer Engineering/Information Technology

Regular Examination May-June 2013

2CE601/2IT601: Theory of Computation

Time: 3 Hours

Total Marks: 70

Instructions:

1. All questions are compulsory
2. Figures to the right indicate full marks.
3. Answer Both Sections in Separate Answer sheets

SECTION - I

Q-1 [A] Give recursive definition of each of followings. [4]

1. The set N of all natural numbers.
2. The language of strings of palindromes.

[B] In each case, say whether the statement is a tautology, a contradiction or neither. Give supporting proof. [4]

1. $Q \wedge (P \rightarrow Q)$
2. $(P \rightarrow Q) \wedge (\neg P \rightarrow Q)$

[C] Prove using mathematical induction that for every $n \geq 1$ [4]

$$\sum i \cdot 2^i = (n-1) \cdot 2^{n+1} + 2$$

OR

Q-1 [A] Write recursive definition of language of strings with more a's than b's. Prove that $REV(x y) = REV(y) REV(x)$ for two string x, y of some alphabet. [6]

[B] State whether following statements are true or false. Also give reason. [4]

1. if $1+1=3$, then $1+2=3$
2. $2+4=3$ only if $2+2=4$

[C] Show that the statements $P \vee Q \vee R \vee S$ and $(\sim P \wedge \sim Q \wedge \sim R) \rightarrow S$ are equivalent. [2]

Q-2 [A] Find the regular expression for FA in following Fig. A and B. [6]

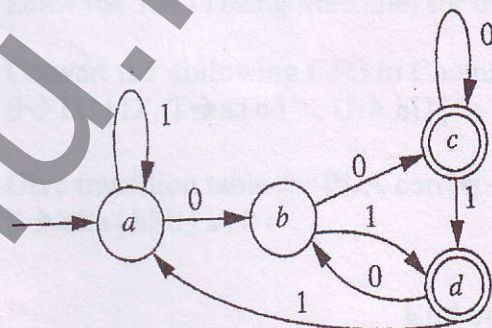


Fig. A.

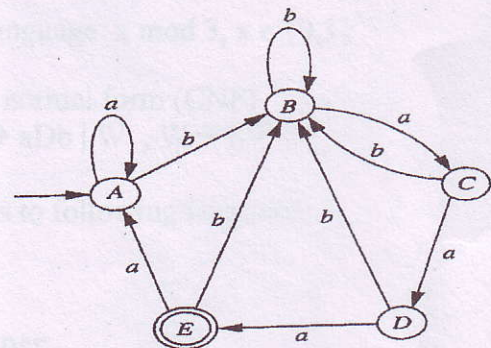


Fig. B.

[B] Convert following NFA given in Fig. B to DFA using subset construction. [5]

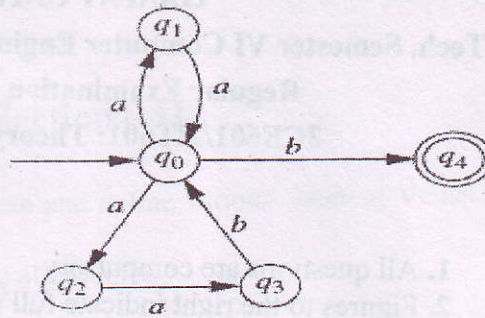


Fig. B.

OR

Q-2 [A] In each case, find a string of minimum length in $\{0,1\}^*$ not in the language corresponding to the given regular expression. [6]

1. $1^*(01)^*0^*$
2. $(0^* + 1^*)(0^* + 1^*)(0^* + 1^*)$
3. $0^*(100^*)1^*$

[B] Draw DFA corresponds to following regular expressions [5]

1. $(11 + 10)^*$
2. $(0+1)(0+1)(0+1)$

Q-3 [A] Draw NFA-null using KLEEN's theorem part-I for $(0+1)^*0^*(10+01)^*$ [6]

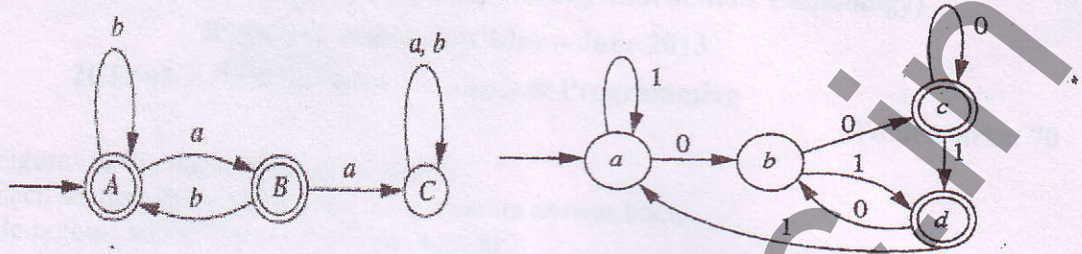
[B] Let $M=(Q,\Sigma,q_0, A, \delta)$ be an NFA - null, where $Q=\{q_0,q_1,q_2,q_3\}$, $\Sigma=\{0,1\}$, $A=\{q_3\}$ and δ is given by the following table. [6]

q	$\delta(q,0)$	$\delta(q,1)$	$\delta(q,\wedge)$
q_0	$\{q_0, q_1\}$	$\{q_0\}$	$\{q_1\}$
q_1	$\{q_2\}$	$\{q_1\}$	$\{q_0, q_2\}$
q_2	$\{q_3\}$	$\{q_3\}$	$\{q_3\}$
q_3	\emptyset	\emptyset	$\{q_1\}$

find $\delta^*(q_0,11)$, $\delta^*(q_0,001)$

SECTION - II

Q-4 [A] Derive regular CFG for the following FA. [6]



[B] Using pumping lemma for regular language prove that language $L = \{0^i 1^i \mid i \geq 0\}$ is not a regular language. [6]

OR

Q-4 [A] Give a grammar which has "Dangling Else" problem. Also draw the derivation tree. [6]

[B] Show, using the pumping lemma for CFL that the given language L is not a CFL. [6]
Where, $L = \{a^n b^{2n} a^n \mid n \geq 0\}$

Q-5 [A] Draw a TM (Turing Machine) accepting language of palindrome strings. [6]

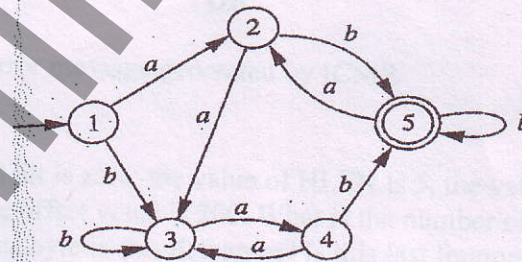
[B] Consider the language of all balanced strings involving two types of brackets $[] \{\}$. [5]

CFG with productions are :

$S \rightarrow SS \mid [S] \mid \Lambda$
Give transition table for **Deterministic PDA** accepting the above language.

OR

Q-5 [A] For the following FA find a minimum state FA recognizing the same language. [6]



[B] Draw the TM (Turing Machine) for the language $x \bmod 3, x \in \{0,1\}^*$ [5]

Q-6 [A] Convert the following CFG to Chomsky normal form (CNF) [6]

$S \rightarrow TU \mid D, T \rightarrow aTb \mid \Lambda, U \rightarrow bU \mid \Lambda, D \rightarrow aDb \mid W, W \rightarrow bW \mid \Lambda$

[B] Give transition table for PDA corresponds to following language [6]

$S \rightarrow aSa \mid bSb \mid a \mid b \mid \Lambda$

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